

Week 13: Secret Sharing and Threshold Cryptography

Secret Sharing

➤ Background

- ✓ Some secrets are too important to be kept by one person.
- ✓ *"It is easier to trust the many than the few"*
- ✓ Secrecy (trust) and robustness

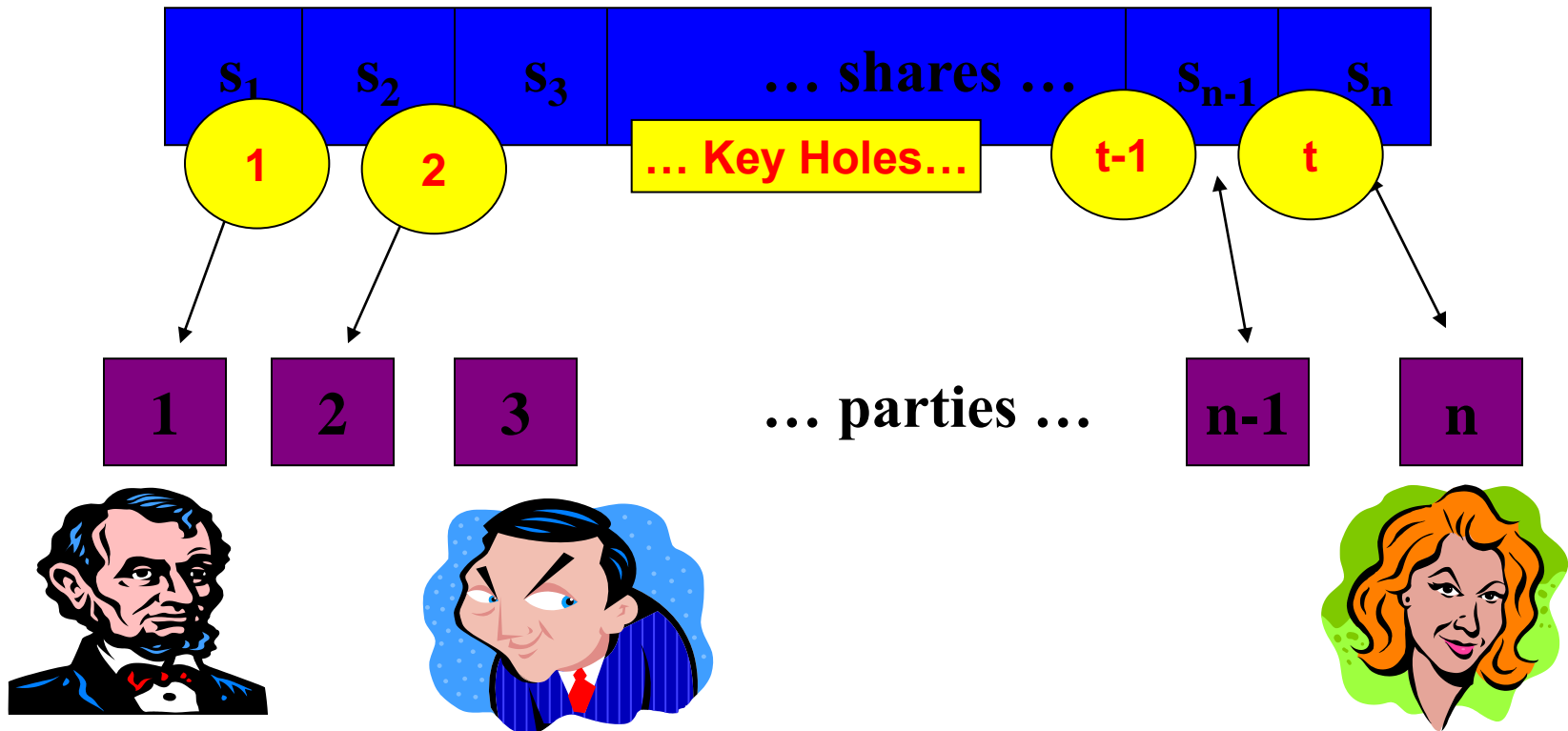
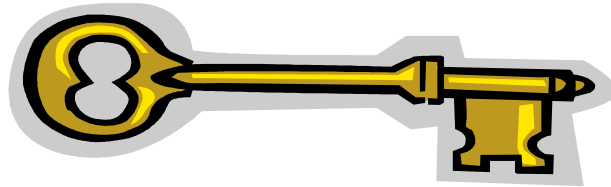
➤ Example:

- ✓ Purported by Time Magazine in 1992 that the Russian nuclear weapon systems were protected by a two-out-of-three access mechanism – President, Defense Minister and Defense Ministry

➤ Secret Sharing

- ✓ Distribute a secret amongst a group of participants
- ✓ Each participant is allocated a share of the secret
- ✓ Secret can be reconstructed only when the shares are combined together
- ✓ Individual shares are of no use on their own.

Secret Sharing (Schematic)



Secret Sharing

➤ Flawed secret sharing



➤ Trivial secret sharing

A secret s is distributed as $s = b_1 \oplus b_2 \oplus \dots \oplus b_{n-1} \oplus b_n$

1) Choose random numbers b_1, \dots, b_{n-1}

2) Compute $b_n = b_1 \oplus b_2 \oplus \dots \oplus b_{n-1} \oplus s$

**All n shares should be present to recover the secret s
(Not robust)**

Wrong Secret Sharing

➤ Flawed secret sharing

password → pa ss wo rd



➤ Trivial secret sharing

A secret s is distributed as $s = b_1 \oplus b_2 \oplus \dots \oplus b_{n-1} \oplus b_n$

- 1) Choose random numbers b_1, \dots, b_{n-1}
- 2) Compute $b_n = b_1 \oplus b_2 \oplus \dots \oplus b_{n-1} \oplus s$

All n shares should be present to recover the secret s
(Not robust)

Threshold Secret Sharing

➤ Scenario

- For example, imagine that the Board of Directors of Coca-Cola would like to protect **Coke's secret formula**. The president of the company should be able to access the formula when needed, but in an emergency any **3 of the 12 board members** would be able to unlock the secret formula together.
- This can be accomplished by a secret sharing scheme with $t = 3$ and $n = 15$, where **3 shares are given to the president, and 1 is given to each board member.**

➤ Security Issues

- **Secrecy**: resistance against any misbehavior
- **Robustness**: reliability against any possible error

SS by Shamir(1/3)

➤ (t, n) Secret Sharing

✓ Secret information K

✓ n share holders (P_1, \dots, P_n)

✓ Using $t-1$ degree random polynomial with random coefficient

(Step 1. Polynomial construction) A dealer selects a secret, K ($< p$: prime) as a constant term and $t-1$ degree random polynomial with arbitrary coefficients as :

$$F(x) = K + a_1x + a_2x^2 + \dots + a_{t-1}x^{t-1} \text{ mod } p$$

(Step 2. Share distribution) Distributes $F(i)$ ($i=1, \dots, n$) securely to share holders P_i .

(Step 3. Secret recovery) When t shares $\Lambda=(K_1, K_2, \dots, K_t)$ among n are given, recover K by using the Lagrange Interpolation

$$K = \sum_{j \in \Lambda} K_j \lambda_{j, \Lambda} \text{ mod } p, \quad \text{where } \lambda_{j, \Lambda} = \prod_{l \in \Lambda \setminus \{j\}} \frac{l}{l - j}$$

SS by Shamir(2/3)

✓ Setup

✓ (3,5) secret sharing

✓ **K=11**, p=17

✓ Construct a degree 2 random polynomial

$$F(x) = K + a_1x + a_2x^2 \text{ mod } p$$

✓ For a random choice $a_1=8$, $a_2=7$

$$F(x) = 11 + 8x + 7x^2 \text{ mod } 17$$

✓ GENSHARE

✓ Share distribution

$$K_1 = F(1) = 7 \times 1^2 + 8 \times 1 + 11 \equiv 9 \pmod{17}$$

$$K_2 = F(2) = 7 \times 2^2 + 8 \times 2 + 11 \equiv 4 \pmod{17}$$

$$K_3 = F(3) = 7 \times 3^2 + 8 \times 3 + 11 \equiv 13 \pmod{17}$$

$$K_4 = F(4) = 7 \times 4^2 + 8 \times 4 + 11 \equiv 2 \pmod{17}$$

$$K_5 = F(5) = 7 \times 5^2 + 8 \times 5 + 11 \equiv 5 \pmod{17}$$

K_1, K_2, K_3, K_4, K_5 : shares given to (P_1, \dots, P_5)

SS by Shamir (3/3)

Using the Lagrange interpolation

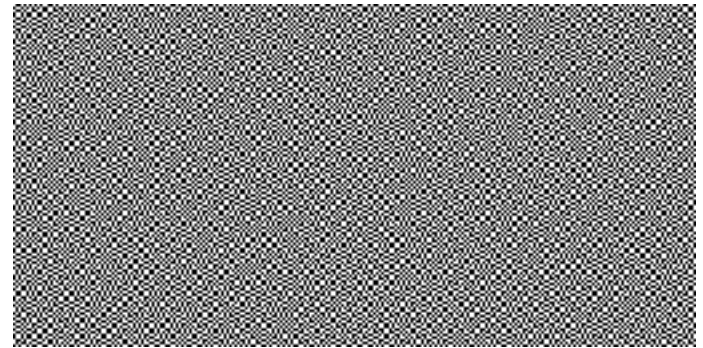
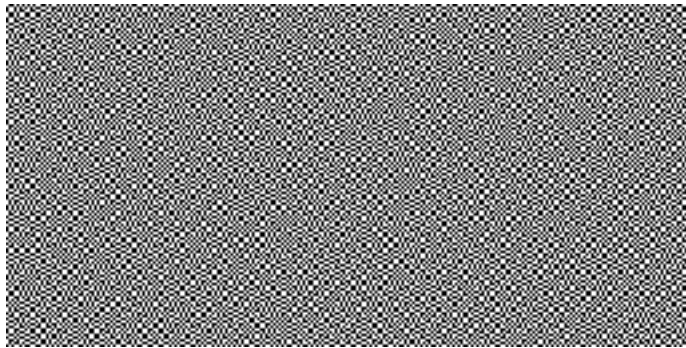
For $\Lambda=(K_1, K_2, K_3)$

$$K = K_1 \left(\frac{2}{2-1} \frac{3}{3-1} \right) + K_2 \left(\frac{1}{1-2} \frac{3}{3-2} \right) + K_3 \left(\frac{1}{1-3} \frac{2}{2-3} \right)$$
$$= 9 \cdot 3 + 4 \cdot (-3) + 13 \cdot 1 \pmod{17} = 11$$

(Quiz) Using $\Lambda=(K_2, K_4, K_5)$, recover secret, **K**

(B&W) Visual Cryptography

- What?



- It is different from the concept of traditional cryptography
- It depends on perception by the human eyes

Color Visual Cryptography

